# EMULATION, REDUCTION, AND EMERGENCE IN DYNAMICAL SYSTEMS Marco Giunti

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#### Résumé

Emergence et réduction sont traditionnellement considérées des catégories incompatibles. Dans cet article je montre que, contrairement à cette idée, émergence et réduction peuvent coexister. Pour étayer cette thèse, je considère les systèmes dynamiques et, sur la base d'un théorème général de représentation, je montre que, pour ces systèmes, la relation d'émulation est suffisante pour la réduction (intuitivement, un système dynamique  $DS_1$  émule un deuxième système dynamique  $DS_2$  quand  $DS_1$  reproduit exactement la dynamique de  $DS_2$ ). Cette vue représentationnelle de la réduction, contrairement à la vue déductiviste traditionnelle, est compatible avec l'existence de propriétés structurelles du système réduit qui ne sont pas aussi des propriétés du système réducteur. Ainsi, de ce point de vue, réduction et émergence ne sont pas du tout des catégories incompatibles mais plutôt complémentaires.

#### **Abstract**

The received view about emergence and reduction is that they are incompatible categories. I argue in this paper that, contrary to the received view, emergence and reduction can hold together. To support this thesis, I focus attention on dynamical systems and, on the basis of a general representation theorem, I argue that, as far as these systems are concerned, the emulation relationship is sufficient for reduction (intuitively, a dynamical system  $DS_1$  emulates a second dynamical system  $DS_2$  when  $DS_1$  exactly reproduces the whole dynamics of  $DS_2$ ). This representational view of reduction, contrary to the standard deductivist one, is compatible with the existence of structural properties of the reduced system that are not also properties of the reducing one. Therefore, under this view, by no means are reduction and emergence incompatible categories but, rather, complementary ones.

#### 1. Introduction

Emergence and reduction are traditionally viewed as incompatible categories (Beckermann 1992<sup>1</sup>; Kim 1992<sup>2</sup>). A property of a high level system is said to be emergent if it cannot be explained in terms of properties of the system's constitutive parts or, more precisely, if it is not one of the properties of more basic parts, which, together, make up the system. Thus, in order to speak of an emergent property P of system  $S_2$  we need to verify, first, that  $S_2$  is made up of another system  $S_1$  (intuitively,  $S_1$  is the system of the constitutive parts of  $S_2$ 

Beckermann, Ansgar (1992), "Supervenience, Emergence and Reduction", in Ansgar Beckermann, Tommaso Toffoli, and Jaegwon Kim (eds.), Emergence or Reduction? Essays on the Prospects of Nonreductive Physicalism. Berlin: Walter de Gruyter, 94-118.

<sup>&</sup>lt;sup>2</sup> Kim, Jaegwon (1992), "Downward Causation in Emergentism and Non-reductive Physicalism", in Ansgar Beckermann, Tommaso Toffoli, and Jaegwon Kim (eds.), *Emergence or Reduction? Essays on the Prospects of Nonreductive Physicalism*. Berlin: Walter de Gruyter, 119-138.

taken in isolation, or in relations different from those typical of  $S_2$ ; see Broad 1925<sup>3</sup>) and, second, that P is not one of the properties of  $S_1$ . But then, the concept of emergence seems to yield a paradox: On the one hand, since  $S_2$  is made up of  $S_1$ ,  $S_2$  is reduced to  $S_1$ ; on the other one, since the property P of  $S_2$  is not one of the properties of  $S_1$ ,  $S_2$  is not reduced to  $S_1$ . The traditional solution denies that the constitution relationship ( $S_2$ 's being made up of  $S_1$ ) is sufficient for reduction. By contrast, the second horn of the dilemma is not questioned, for it is taken for granted that  $S_2$ 's reduction to  $S_1$  entails that any property of  $S_2$  is also a property of  $S_1$ .

This paper maintains that the traditional solution is irremediably flawed. In fact, there are pairs of systems,  $S_2$  and  $S_1$ , for which both the constitution relationship ( $S_2$  is made up of  $S_1$ ) and the reduction one ( $S_2$  is reduced to  $S_1$ ) clearly hold together. Moreover, for these pairs of systems, it also turns out that some property of  $S_2$  is not a property of  $S_1$ , so that any such property is emergent. It follows that, contrary to the received view, emergence and reduction by no means are incompatible categories but, rather, complementary ones.

To support this thesis, I will consider some simple examples of dynamical systems for which the emulation relationship holds. As intended here (Arnold 1977<sup>4</sup>; Szlensk 1984<sup>5</sup>; Giunti 1997<sup>6</sup>), a dynamical system is a mathematical model that expresses the idea of an arbitrary deterministic system, either reversible or irreversible, with discrete or continuous time or state space. Such models allow us to study in a precise way a whole series of typical phenomena in complex systems. Among them, in recent years, the phenomenon of emulation has gained growing attention (Wolfram 1983a<sup>7</sup>, 1983b<sup>8</sup>, 1984a<sup>9</sup>, 1984b<sup>10</sup>,  $2002^{11}$ ). Intuitively, a dynamical system  $DS_1$  emulates a second dynamical system  $DS_2$ when the first one exactly reproduces the whole dynamics of the second one. emulation relationship can be defined in a precise way for any two arbitrary dynamical systems, and it has also been shown (Giunti 1997<sup>12</sup>, ch.1, th. 11) that, if  $DS_1$  emulates  $DS_2$ , there is a third system  $DS_3$  such that (i)  $DS_2$  is isomorphic to  $DS_3$ ; (ii) all states of  $DS_3$  are states of  $DS_1$ ; (iii) any state transition of  $DS_3$  is constructed out of state transitions of  $DS_1$ . Because of this result, it makes perfect sense to claim that  $DS_2$  is made up of  $DS_1$ , as well as that  $DS_2$  is reduced to  $DS_1$ . Therefore, to show that both reduction and emergence can hold together, it suffice to exhibit two dynamical systems  $DS_1$  and  $DS_2$ , as well as a property P, such that  $DS_1$  emulates  $DS_2$ ,  $DS_2$  has P, but  $DS_1$  does not have P. I will show that this

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Broad, Charlie Dunbar (1925), *The Mind and its Place in Nature*. London: Routledge and Kegan Paul.

Arnold, Vladimir I. (1977), Ordinary Differential Equations. Cambridge: The MIT Press.
Szlensk, Wieslaw (1984), An Introduction to the Theory of Smooth Dynamical Systems. Chichister, England: John Wiley and Sons.

<sup>&</sup>lt;sup>6</sup> Giunti, Marco (1997), Computation, Dynamics, and Cognition. New York: Oxford University Press.

Wolfram, Stephen (1983a), "Statistical Mechanics of Cellular Automata", Reviews of Modern Physics 55, 3:601-644.

<sup>&</sup>lt;sup>8</sup> Wolfram, Stephen (1983b), "Cellular Automata", Los Alamos Science 9:2-21.

Wolfram, Stephen (1984a), "Computer Software in Science and Mathematics", Scientific American 56:188-203.

Wolfram, Stephen (1984b), "Universality and Complexity in Cellular Automata", in Doyne Farmer, Tommaso Toffoli, and Stephen Wolfram (eds.), Cellular Automata. Amsterdam: North Holland Publishing Company, 1-35.

<sup>&</sup>lt;sup>11</sup> Wolfram, Stephen (2002), A New Kind of Science. Champaign. IL: Wolfram Media, Inc.

<sup>&</sup>lt;sup>12</sup> See note 6.

situation already obtains for two pairs of simple finite discrete systems and that, in either case, the emergent property P is a strong form of irreversibility of system  $DS_2$ .

### 2. Dynamical systems and emulation

A dynamical system is a mathematical model that expresses the idea of an arbitrary deterministic system, either reversible or irreversible, with discrete or continuous time or state space. Let Z be the integers,  $Z^+$  the non-negative integers, R the reals and  $R^+$  the non-negative reals; below is the exact definition of a dynamical system.

- [1] DS is a dynamical system iff there is M, T,  $(g^t)_{t \in T}$  such that  $DS = (M, (g^t)_{t \in T})$  and
  - 1. *M* is a non-empty set; *M* represents all the possible states of the system, and it is called the *state space*;
  - 2. T is either Z,  $Z^+$ , R, or  $R^+$ ; T represents the time of the system, and it is called the *time set*;
  - 3.  $(g^t)_{t \in T}$  is a family of functions from M to M; each function  $g^t$  is called a *state transition* or a *t-advance* of the system;
  - 4. for any  $t, v \in T$ , for any  $x \in M$ ,  $g^{0}(x) = x$  and  $g^{t+v}(x) = g^{v}(g^{t}(x))$ .
- [2] A discrete dynamical system is a dynamical system whose state space is finite or denumerable, and whose time set is either Z or  $Z^+$ ; examples of discrete dynamical systems are Turing machines and cellular automata. [3] A continuous dynamical system is a dynamical system that is not discrete; examples of continuous dynamical systems are iterated mappings on R, and systems specified by ordinary differential equations.
- [4]  $DS = (M, (g^t)_{t \in T})$  is a possible dynamical system iff DS satisfies the first three conditions of definition [1]. We can now define the concept of an isomorphism between two possible dynamical systems as follows. [5] u is an isomorphism of  $DS_1$  in  $DS_2$  iff  $DS_1 = (M, (g^t)_{t \in T})$  and  $DS_2 = (N, (h^v)_{v \in V})$  are possible dynamical systems, T = V,  $u: M \to N$  is a bijection and, for any  $t \in T$ , for any  $x \in M$ ,  $u(g^t(x)) = h^t(u(x))$ .
- [6]  $DS_1$  is isomorphic to  $DS_2$  iff there is u such that u is an isomorphism of  $DS_1$  in  $DS_2$ . It is easy to verify that the isomorphism relation is an equivalence relation on any given set of possible dynamical systems. (The concept of set of all possible dynamical systems is inconsistent, and we must then take as the basis of the theory of dynamical systems a specific, sufficiently large, set of possible dynamical systems.)

It is also not difficult to prove that the relation of isomorphism is a congruence with respect to the property of being a dynamical system, that is to say: if  $DS_1$  is isomorphic to  $DS_2$  and  $DS_1$  is a dynamical system, then  $DS_2$  is a dynamical system. This allows us to speak of abstract dynamical systems in exactly the same sense we talk of abstract groups, fields, lattices, order structures, etc. We can thus define: [7] an abstract dynamical system is any equivalence class of isomorphic dynamical systems.

Dynamical systems are appropriate models to study several interesting phenomena in complex systems. The one of emulation is typical of computational systems (Wolfram 2002<sup>13</sup>), but it can in principle involve any two dynamical systems. The intuitive idea is

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<sup>&</sup>lt;sup>13</sup> See note 11.

that a dynamical system  $DS_1$  emulates a second dynamical system  $DS_2$  when the first one exactly reproduces the whole dynamics of the second one. Here are some examples. A universal Turing machine emulates any Turing machine; for any Turing machine TM there is a cellular automaton CA such that CA emulates TM (Smith 1971<sup>14</sup>, th. 3), and vice versa; the simple cellular automaton specified by Wolfram's rule 18 emulates the one specified by rule 90 (both CA are monodimensional, with 2 possible values for cell, and neighborhood of radius 1; see Wolfram 1983b<sup>15</sup>, 20).

Giunti 1997<sup>16</sup> (ch. 1, def. 4) gave a formal definition of the emulation relationship that applies to any two arbitrary dynamical systems. Here, I will employ a weaker and simpler definition, which nevertheless suffices for the present purposes.

[8]  $DS_1$  emulates  $DS_2$  iff  $DS_1 = (M, (g^t)_{t \in T})$  and  $DS_2 = (N, (h^v)_{v \in V})$  are dynamical systems, and there is an injective function  $u: N \to M$  such that, for any  $c \in N$ , for any  $v \in V$ , there is  $t \in T$  such that  $u(h^v(c)) = g^t(u(c))$ . Any function u that satisfies the previous condition is called an emulation of  $DS_2$  in  $DS_1$ .

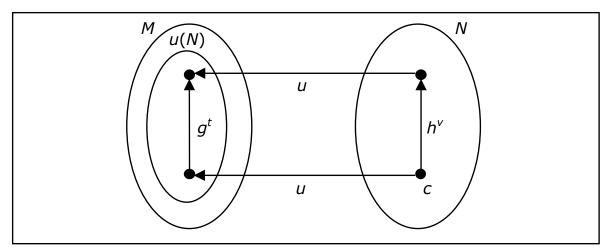


FIGURE 1 Emulation

## 3. Emulation, constitution, and reduction

Giunti 1997<sup>17</sup> (ch.1, th. 11) proved that, if u is an emulation of  $DS_2$  in  $DS_1$ , there is a third system  $DS_3$  such that (i) u is an isomorphism of  $DS_2$  in  $DS_3$ ; (ii) all states of  $DS_3$  are states of  $DS_1$ ; (iii) any state transition of  $DS_3$  is constructed out of state transitions of  $DS_1$ . This result still holds for the weaker definition of emulation [8], as the following theorem shows.

<sup>&</sup>lt;sup>14</sup> Smith, Alvy Ray III (1971), "Simple Computation-universal Cellular Spaces", *Journal of the Association for Computing Machinery* 18, 3:339-353.

<sup>15</sup> See note 8.

<sup>&</sup>lt;sup>16</sup> See note 6.

<sup>&</sup>lt;sup>17</sup> See note 6.

Virtual System Theorem [VST]

- Let  $DS_1 = (M, (g^t)_{t \in T})$  and  $DS_2 = (N, (h^v)_{v \in V})$  be dynamical systems, and u be an emulation of  $DS_2$  in  $DS_1$ ;
- let  $DS_3 = (\underline{N}, (\underline{h}^v)_{v \in V})$ , where  $\underline{N} = u(N)$  and, for any  $a \in \underline{N}$ , for any  $v \in V$ ,  $\underline{h}^v(a) = u(h^v(u^{-1}(a)))$ ; the system  $DS_3$  is called *the virtual u-system*  $DS_2$  *in*  $DS_1$  (see figure 2); then:
- (i) u is an isomorphism of  $DS_2$  in  $DS_3$ ;
- (ii) all states of  $DS_3$  are states of  $DS_1$ ;
- (iii) for any state transition  $\underline{h}^{v}$  of  $DS_3$ , for any  $a \in \underline{N}$ , there is a state transition  $g^{t}$  of  $DS_1$  such that  $\underline{h}^{v}(a) = g^{t}(a)$ .

## *Proof of* (i)

By the definition of  $DS_3$ , for any  $c \in N$ ,  $u(h^v(c)) = u(h^v(u^{-1}(u(c)))) = \underline{h}^v(u(c))$ . Therefore, by the definition of isomorphism [5], u is an isomorphism of  $DS_2$  in  $DS_3$ .

### Proof of (ii)

Obvious, by the definition of  $DS_3$ .

## Proof of (iii)

By the definition of  $DS_3$ , for any  $v \in V$ , for any  $a \in \underline{N}$ ,  $\underline{h}^v(a) = u(h^v(u^{-1}(a)))$ . Let  $c = u^{-1}(a)$ . Since u is an emulation of  $DS_2$  in  $DS_1$ , by definition [8], there is  $t \in T$  such that  $u(h^v(c)) = g^t(u(c))$ . Therefore,  $\underline{h}^v(a) = g^t(u(c)) = g^t(a)$ . Q.E.D.

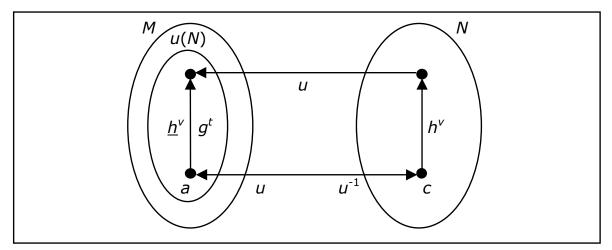


FIGURE 2

The virtual u-system  $DS_2$  in  $DS_1$ 

Because of [VST], if a dynamical system  $DS_1$  emulates a second system  $DS_2$ , it makes perfect sense to claim that  $DS_2$  is made up of  $DS_1$ , as well as that  $DS_2$  is reduced to  $DS_1$ . In other words, I maintain that, in virtue of [VST], emulation is sufficient for both constitution and reduction.

### 4. Emergence and reduction

A property P of a high level system  $S_2$  is said to be *emergent with respect to a lower level system*  $S_1$  just in case (a)  $S_2$  is made up of  $S_1$  (intuitively,  $S_1$  is the system of the constitutive parts of  $S_2$  taken in isolation, or in relations different from those typical of  $S_2$ ; see Broad 1925<sup>18</sup>) and (b) P is not one of the properties of  $S_1$ .

Therefore, since emulation is sufficient for both constitution and reduction, in order to show that emergence and reduction can hold together, it is sufficient to exhibit a pair of dynamical systems  $DS_1$  and  $DS_2$ , as well as a property P, such that  $DS_1$  emulates  $DS_2$ ,  $DS_2$  has P, but  $DS_1$  does not have P. In the next section, I will give two examples of such pairs of systems. For each pair, both  $DS_1$  and  $DS_2$  are small finite discrete systems (with just three states), while the emergent property P is the strong irreversibility  $^{20}$  of system  $DS_2$ .

## 5. Examples of dynamical systems $DS_1$ and $DS_2$ such that (i) $DS_2$ is reduced to $DS_1$ and (ii) $DS_2$ has emergent properties with respect to $DS_1$

To state the examples, we first need a few more general concepts of dynamical systems theory. [9] A *cascade* is a dynamical system with discrete time, i.e., whose time set is either Z or  $Z^+$ . [10] A dynamical system is *reversible* iff its time set is either Z or R; [11] it is *irreversible* iff its time set is either  $Z^+$  or  $R^+$ . Note that any t-advance  $g^t$  (t > 0) of an irreversible cascade (M, ( $g^t$ ) $_{t \in Z^+}$ ) can always be thought as being generated by iterating t times a given function  $g: M \to M$  (thus,  $g^1 = g$ ). Therefore, as far as an irreversible cascade is concerned, the whole dynamics of the system reduces to the behavior of its first t-advance  $g^1$ .

[12] A dynamical system is *logically reversible* iff it is irreversible, but all its state-transitions are injective; [13] it is *logically irreversible* iff it is irreversible and at least one of its state-transitions is not injective; [14] it is *strongly irreversible* iff there are two different states a and b and a state-transition  $g^v$  such that  $g^v(a) = g^v(b)$  and, for any state-transition  $g^t$ ,  $g^t(a) \neq b$  and  $g^t(b) \neq a$ . Obviously, by definitions [12], [13] and [14], if a dynamical system is logically reversible, it is not strongly irreversible.

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<sup>&</sup>lt;sup>18</sup> See note 3.

In order to avoid trivial cases, it is also intended that P be a *structural property* of the kind of structure that both  $S_1$  and  $S_2$  share. This means the following. (i) The two systems  $S_1$  and  $S_2$  are systems of the same mathematical kind K (for example, they are both dynamical systems, or groups, rings, etc.); (ii) the appropriate isomorphism relationship  $\equiv$  is defined for the kind of system K; (iii) the property P is preserved by the isomorphism  $\equiv$ , that is to say, for any two systems  $S_1$  and  $S_2 \in K$ , if  $S_1$  has P and  $S_1 \equiv S_2$ , then  $S_2$  has P; (iv) the property P is specific to the kind of structure K, that is to say, for any system S, if  $S \notin K$ , then S has not P.

<sup>&</sup>lt;sup>20</sup> Strong irreversibility is defined in the next section. It is easy to verify that strong irreversibility is a structural property (see note 19) of dynamical systems.

Figure 3 shows a pair of cascades  $DS_1 = (M, (g^t)_{t \in Z^+})$  and  $DS_2 = (N, (h^v)_{v \in Z^+})$  such that (i)  $DS_2$  is reduced to  $DS_1$  and (ii) the property P of strong irreversibility is an emergent property of  $DS_2$  with respect to  $DS_1$ .

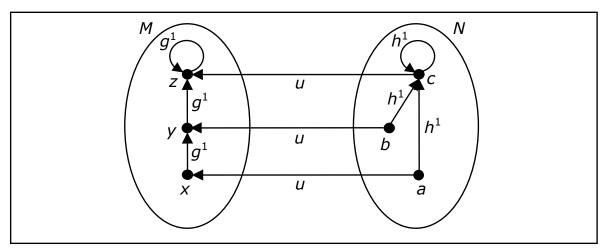


FIGURE 3  $DS_1$  emulates  $DS_2$ ,  $DS_1$  is logically irreversible but not strongly irreversible, and  $DS_2$  is strongly irreversible

Figure 4 shows a second pair of cascades  $DS_1 = (M, (g^t)_{t \in Z^+})$  and  $DS_2 = (N, (h^v)_{v \in Z^+})$  such that (i)  $DS_2$  is reduced to  $DS_1$  and (ii) the property P of strong irreversibility is an emergent property of  $DS_2$  with respect to  $DS_1$ .

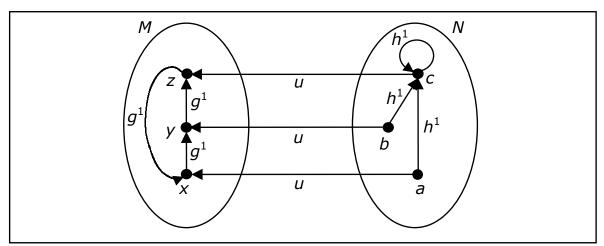


FIGURE 4  $DS_1$  emulates  $DS_2$ ,  $DS_1$  is logically reversible (thus, not strongly irreversible), and  $DS_2$  is strongly irreversible

## 6. Concluding remarks: Toward a general representational theory of reduction and emergence

Traditionally, reduction has been analyzed in terms of a *deductive* relationship between two empirically interpreted formal *theories*, via correspondence rules between the terms of the

two theories (Nagel  $1961^{21}$ ; Churchland  $1979^{22}$ ,  $1985^{23}$ ; Hooker  $1981^{24}$ ). By shifting attention from formal theories to mathematical *models*, it is natural to think of reduction in terms of some kind of *representation* relationship between two models. This paper has argued that, if the two models are dynamical systems, the relationship of emulation is sufficient for reduction (in virtue of [*VST*]).

An important point needs to be stressed. If we think of  $S_2$ 's reduction to  $S_1$  as a form of deduction of theory  $S_2$  from theory  $S_1$  (more precisely, the deduction of a relevantly isomorphic image of  $S_2$  from  $S_1$ ; see Churchland  $1985^{25}$ , sec. 1; Beckermann  $1992^{26}$ , 108), then it is obvious that all the properties of  $S_2$  (more precisely, the properties referred to by the relevantly isomorphic image of  $S_2$ ) are a fortiori properties of  $S_1$ . Therefore, if we take this kind of approach to reduction, there cannot be two theories  $S_2$  and  $S_1$  such that  $S_2$  is reduced to  $S_1$  and  $S_2$  has emergent properties with respect to  $S_1$ .

But this need not be the case if we think of reduction as a form of *representation* between two models  $S_1$  and  $S_2$ , which grants the construction, within the representing model  $S_1$ , of an isomorphic (or, perhaps, just homomorphic) image of  $S_2$ . In fact, as I have just shown for the special case of dynamical systems, this view of reduction is compatible with the existence of structural properties of the reduced system that are not also properties of the reducing one. Therefore, under this view, reduction and emergence no longer are incompatible relationships but, rather, complementary ones.

At present, the *representational theory* of reduction and emergence has a precise formulation only if the models involved are dynamical systems. Even though many interesting models in real science are of this kind, by no means is this special formulation sufficient to account for all relevant cases of reduction or emergence. What we need is a *general* representational theory, as precise as the one restricted to dynamical systems, which apply to *arbitrary models*. The formulation of such a general theory, however, is not an easy matter, for it involves a preliminary investigation of fairly hard questions like: What is, *in general*, a mathematical structure? What is, *in general*, a mathematical model? What is an isomorphism between two *arbitrary* models? What is the relationship between two *arbitrary* models that generalizes the one of emulation between dynamical systems?

<sup>21</sup> Nagel, Ernest (1961), *The Structure of Science*. New York: Harcourt, Brace & World.

<sup>&</sup>lt;sup>22</sup> Churchland, Paul M. (1979), Scientific Realism and the Plasticity of Mind. Cambridge: Cambridge University Press.

<sup>&</sup>lt;sup>23</sup> Churchland, Paul M. (1985), "Reduction, Qualia, and the Direct Introspection of Brain States", *Journal of Philosophy* 82, 1:8-28.

<sup>&</sup>lt;sup>24</sup> Hooker, Clifford Alan (1981), "Towards a General Theory of Reduction", *Dialogue* 20:38-60, 201-236, 496-529.

<sup>&</sup>lt;sup>25</sup> See note 23.

<sup>&</sup>lt;sup>26</sup> See note 1.